

# Speedable left-c.e. numbers

Wolfgang Merkle and Ivan Titov

Universität Heidelberg, Heidelberg, Germany

- 1 Introduction
  - Effective Reals
  - Notion of Randomness
- 2 Independence from the constant and approximation
- 3 Speedable reals
  - Nonspeedability of Martin-Löf random left-c.e. reals
  - Speedable left-c.e. reals and Solovay degrees
- 4 Appendix: generalization of speedability on  $\Delta_0^2$ -reals

- We start by reviewing some notions of effective approximability and of randomness for real numbers.

## Definition

A **computable approximation** of a real  $\alpha$  is a computable sequence  $(a_n)_{n \in \omega}$  of dyadic rationals such that  $\lim_{n \rightarrow \infty} a_n = \alpha$ .

A **left-c.e. approximation** of a real  $\alpha$  is a strictly increasing computable approximation and a **right-c.e. approximation** of a real  $\alpha$  is a strictly decreasing computable approximation.

A real  $\alpha$  is a  $\Delta_0^2$  **real** if there exists a computable approximation of  $\alpha$ .

A real  $\alpha$  is **left-c.e.** if there is a left-c.e. approximation of  $\alpha$  and is **right-c.e.** if there is a right-c.e. approximation of  $\alpha$ .

# Notions of randomness

- A **Martin-Löf test** is a uniformly effective sequence of open sets  $U_0, U_1, \dots$ , such that the set  $U_n$  has uniform measure of at most  $2^{-n-1}$ .
- A **Solovay test** is a uniformly effective sequence of open sets  $S_0, S_1, \dots$ , such that the sum of uniform measures of all  $S_n$  is finite.

## Definition

A sequence  $A$  is **Martin-Löf random**, if there is no Martin-Löf test  $U_0, U_1, \dots$ , such that  $A \in \bigcap_{i \in \omega} U_i$ .

- It is well-known that a sequence  $A$  is Martin-Löf random if and only if there is no Solovay test  $S_0, S_1, \dots$  such that  $A$  is contained in infinitely many of the sets  $S_n$ .

Further, we identify any real  $\alpha := 0.A$  with its binary representation  $A$ .

## Definition

A left-c.e. (or right-c.e.) real  $\alpha$  is  **$\rho$ -speedable with respect to its left-c.e. (or right-c.e.) approximation**  $(a_n)_{n \rightarrow \infty}$  for a constant  $\rho \in (0, 1)$  if there is a computable strictly increasing function  $f : \mathbb{N} \rightarrow \mathbb{N}$  such that

$$\liminf_{n \rightarrow \infty} \frac{\alpha - a_{f(n)}}{\alpha - a_n} < \rho.$$

Such a function  $f$  is called a **speed-up function**.

A left-c.e. or right-c.e. real  $\alpha$  is  **$\rho$ -speedable** if it is  $\rho$ -speedable with respect to some of its left-c.e. (or right-c.e.) approximations.

- When replacing  $\liminf$  by  $\limsup$  in the definition above, the newly defined concept holds exactly for every computable real  $\alpha$ .

# Independence of speedability from the constant and approximation

## Theorem

*A left-c.e. real  $\alpha$  that is  $\rho$ -speedable for some  $\rho \in (0, 1)$  is  $\rho'$ -speedable with respect to all of its left-c.e. approximations and for all  $\rho' \in (0, 1)$ .*

- A similar statement holds for right-c.e. reals.
- The theorem allows us to use the notion of speedability on the field of left-c.e. (or right-c.e.) reals without specifying a left-c.e. (or right-c.e.) approximation and a speeding constant.

# Independence from the constant and approximation choice: Proof sketch

The proof follows from three statements:

## Lemma

*If  $\alpha$  is  $\rho$ -speedable with respect to some of its left-c.e. approximations  $(a_n)_{n \rightarrow \infty}$ , then  $\alpha$  is  $\rho$ -speedable with respect to  $(a_{f^{[2n]}(0)})_{n \rightarrow \infty}$  or  $(a_{f^{[2n+1]}(0)})_{n \rightarrow \infty}$  via the speed-up function  $n \mapsto n + 1$ .*

## Lemma

*If  $\alpha$  is  $\rho$ -speedable with respect to some of its left-c.e. approximations  $(a_n)_{n \rightarrow \infty}$ , then  $\alpha$  is  $\rho'$ -speedable for every  $\rho' \in (0, 1)$  with respect to the same approximation.*

## Lemma

*If  $\alpha$  is  $\rho$ -speedable, then  $\alpha$  is  $\rho$ -speedable with respect to every of its left-c.e. approximations.*

## Theorem

*Every Martin-Löf random left-c.e. real is nonspeedable.*

## Proof.

Given a speedable left-c.e. real  $\alpha$ , by the proof of the previous theorem we pick some left-c.e. approximation  $a_0, a_1, \dots$  of  $\alpha$  that is  $1/3$ -speedable via the successor function  $n \mapsto n + 1$ , so

$$(a_0, a_0 + 3(a_1 - a_0)), (a_1, a_1 + 3(a_2 - a_1)), \dots$$

is a Solovay test with sum of measures  $3(\alpha - a_0)$ , which overlaps  $\alpha$  infinitely many times. This contradicts the Martin-Löf randomness of  $\alpha$ . □



# Speedable left-c.e. reals and Solovay degrees

- By the preceding theorem, all left-c.e. reals in the maximum Solovay degree are nonspeedable.
- It can also be shown that all nonhigh left-c.e. reals are speedable.
- Speedability is indeed a degree property for Solovay degrees of left-c.e. reals, i.e., two left-c.e. reals that are both in the same Solovay degree are either both speedable or both nonspeedable.
- It is open whether there are nonspeedable left-c.e. reals outside the maximum Solovay degree, or alternatively, whether there are nonspeedable left-c.e. reals that are not Martin-Löf random.

# Speedability of pure $\Delta_0^2$ -reals

One can extend the notion of speedability to the general case of  $\Delta_0^2$ -reals in different ways depending on the meaning of "capturing"  $\alpha$ . In this part we introduce the most intuitive generalization which consider the reducing of the distance to  $\alpha$ .

## Definition

A  $\Delta_0^2$  real  $\alpha$  is  **$\rho$ -speedable with respect to its computable approximation**  $(a_n)_{n \rightarrow \infty}$  for a constant  $\rho \in (0, 1)$  if there is a computable strictly increasing function  $f : \mathbb{N} \rightarrow \mathbb{N}$  such that

$$\liminf_{n \rightarrow \infty} \frac{|\alpha - a_{f(n)}|}{|\alpha - a_n|} < \rho.$$

# Speedability of pure $\Delta_0^2$ -reals

## Definition

A  $\Delta_0^2$  real is **pure** if it is neither left-c.e. nor right c.e.

## Theorem

*Every pure  $\Delta_0^2$  real is  $\rho$ -speedable for every constant  $\rho \in (0, 1)$  with respect to each of its computable approximations.*

## Proof.

Omitted. □